

Lecture 4: Expectation

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October 22, 2024

Outline

- 1 Expectation & Variance
- 2 Geometric and Negative Binomial
- 3 Indicator R.V.s and The Fundamental Bridge
- 4 Moments and Indicators
- 5 Poisson
- 6 Distance between Two Probability Distributions
- 7 Probability Generating Functions
- 8 Reading for Fun

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Expectation of A Discrete R.V.

Definition

The *expected value* (also called the *expectation* or *mean*) of a discrete r.v. X whose distinct possible values are x_1, x_2, \dots is defined by

$$E(X) = \sum_{j=1}^{\infty} x_j P(X = x_j)$$

If the support is finite, then this is replaced by a finite sum. We can also write

$$E(X) = \sum_x \underbrace{x}_{\text{value}} \underbrace{P(X = x)}_{\text{PMF at } x}$$

where the sum is over the support of X .

Distribution

Theorem

If X and Y are discrete r.v.s with the same distribution, then $E(X) = E(Y)$ (if either side exists).

Linearity

The expected value of a sum of r.v.s is the sum of the individual expected values.

Theorem

For any r.v.s X , Y and any constant c ,

$$E(X + Y) = E(X) + E(Y)$$

$$E(cX) = cE(X)$$

Monotonicity of Expectation

Theorem

Let X and Y be r.v.s such that $X \geq Y$ with probability 1. Then $E(X) \geq E(Y)$, with equality holding if and only if $X = Y$ with probability 1.

Expectation via Survival Function

Theorem

Let X be a nonnegative integer-valued r.v. Let F be the CDF of X , and $G(x) = 1 - F(x) = P(X > x)$. The function G is called the survival function of X . Then

$$E(X) = \sum_{n=0}^{\infty} G(n)$$

That is, we can obtain the expectation of X by summing up the survival function (or, stated otherwise, summing up tail probabilities of the distribution).

Proof

Law Of The Unconscious Statistician (LOTUS)

Theorem

If X is a discrete r.v. and g is a function from \mathbb{R} to \mathbb{R} , then

$$E(g(X)) = \sum_x g(x) P(X = x)$$

where the sum is taken over all possible values of X .

Variance and Standard Deviation

Definition

The variance of an r.v. X is

$$\text{Var}(X) = E(X - EX)^2.$$

The square root of the variance is called the *standard deviation (SD)*:

$$\text{SD}(X) = \sqrt{\text{Var}(X)}.$$

Properties of Variance

- For any r.v. X , $\text{Var}(X) = E(X^2) - (EX)^2$.
- $\text{Var}(X + c) = \text{Var}(X)$ for any constant c .
- $\text{Var}(cX) = c^2 \text{Var}(X)$ for any constant c .
- If X and Y are independent, then $\text{Var}(X + Y) = \text{Var}(X) + \text{Var}(Y)$.
- $\text{Var}(X) \geq 0$ with equality if and only if $P(X = a) = 1$ for some constant a .

Properties of Variance

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Story: Geometric Distribution

Consider a sequence of independent Bernoulli trials, each with the same success probability $p \in (0, 1)$, with trials performed until a success occurs. Let X be the number of **failures** before the first successful trial. Then X has the Geometric distribution with parameter p ; we denote this by $X \sim \text{Geom}(p)$.

Geometric PMF

Theorem

If $X \sim \text{Geom}(p)$, then the PMF of X is

$$P(X = k) = q^k p$$

for $k = 0, 1, 2, \dots$, where $q = 1 - p$.

Memoryless Property

Theorem

If $X \sim \text{Geom}(p)$, then for any positive integer n ,

$$P(X \geq n + k | X \geq k) = P(X \geq n)$$

for $k = 0, 1, 2, \dots$

Memoryless Property

Theorem

Suppose for any positive integer n , discrete random variable X satisfies

$$P(X \geq n + k | X \geq k) = P(X \geq n)$$

for $k = 0, 1, 2, \dots$, then $X \sim \text{Geom}(p)$.

Memoryless Property

Theorem

Geometric distribution is the one and the only one discrete distribution that is memoryless.

First Success Distribution

Definition

In a sequence of independent Bernoulli trials with success probability p , let Y be the number of trials until the first successful trial, including the success. Then Y has the First Success distribution with parameter p ; we denote this by $Y \sim \text{FS}(p)$.

Example: Geometric & First Success Expectation

Let $X \sim \text{Geom}(p)$ and $Y \sim \text{FS}(p)$, find $E(X)$ and $E(Y)$.

Story: Negative Binomial Distribution

In a sequence of independent Bernoulli trials with success probability p , if X is the number of failures before the r^{th} success, then X is said to have the Negative Binomial distribution with parameters r and p , denoted $X \sim \text{NBin}(r, p)$.

Negative Binomial PMF

Theorem

If $X \sim \text{NBin}(r, p)$, then the PMF of X is

$$P(X = n) = \binom{n+r-1}{r-1} p^r q^n$$

for $n = 0, 1, 2, \dots$, where $q = 1 - p$.

Geometric & Negative Binomial

Theorem

Let $X \sim \text{NBin}(r, p)$, viewed as the number of failures before the r th success in a sequence of independent Bernoulli trials with success probability p . Then we can write $X = X_1 + \cdots + X_r$ where the X_i are i.i.d. $\text{Geom}(p)$.

Example: Expectation

Let $X \sim NBin(r, p)$, find $E(X)$.

Example:



Example:



水浒英雄传

行者·武松

天罡：天伤星
职位：步军副大将
武器：青龙白虎双刀
火杀技：天伤日月斩 玉环醉步
破云无敌脚

攻击力：
攻击范围：
防御力：

15
45
83

人物小传
曾在景阳冈醉打猛虎，名扬天下。为替武大郎报仇，怒杀潘金莲，西门庆，被发配孟州。在孟州，为替“金眼彪”施恩夺回“快活林”酒楼，醉打蒋门神。后遭张都监陷害，武松大闹飞云浦，血溅鸳鸯楼。三山聚义后，在梁山坐第十四把交椅，任方腊时损失一臂手臂。此后在六和寺颐养天年，封为清忠祖师。

统一 小浣熊

Example:



水浒传英雄传

豹子头·林冲

天罡：天雄星
职位：五虎大将之右军大将
武器：寒星冷月枪
必杀技：寒星夺魄刺 冷月索命舞
忧那飞花

攻击力：
防御力：
60 20 95

人物小传
东京八十万禁军教头，因为高俅之子高衙内垂涎林冲妻子，高俅伙同陆谦设计使林冲误闯白虎节堂，被发配沧州途中在野猪林由鲁智深搭救，幸免遇难，但高俅又使人火烧草料场，林冲才大闹山神庙，杀陆谦，雪夜上梁山。后来火并王伦，拥立晁盖为梁山之主。梁山座次排在第六，一杆枪，天下无敌，从无数挂。

6 统一 小浣熊

Example:

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Example:



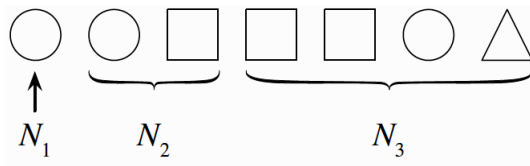
Example:



Model: Coupon Collector

Suppose there are n types of toys, which you are collecting one by one, with the goal of getting a complete set. When collecting toys, the toy types are random (as is sometimes the case, for example, with toys included in cereal boxes or included with kids' meals from a fast food restaurant). Assume that each time you collect a toy, it is equally likely to be any of the n types. Let N denote the number of toys needed until you have a complete set. Find $E(N)$ and $\text{Var}(N)$.

Solution: Coupon Collector



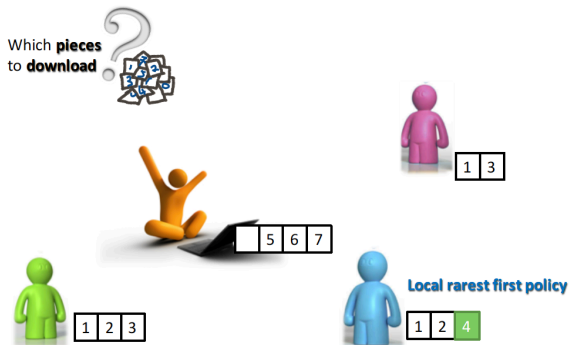
Solution: Coupon Collector

Application: Peer-to-Peer System

- Target file is decomposed into n pieces.
- Each peer randomly downloads pieces and uploads pieces from its neighbors.
- $\Theta(n \ln n)$ downloads to complete the downloading file.
- The last block problem: missing the last piece (stop at 99% downloading progress)

Application: Peer-to-Peer System

- Solution adopted by BitTorrent:
 - ▶ tries to download a block that is least replicated among its neighbors
 - ▶ maximize the diversity of content in the system, i.e., make the number of replicas of each block as equal as possible.



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Properties of Indicator R.V.

Let A and B be events. Then the following properties hold.

① $(I_A)^k = I_A$ for any positive integer k .

② $I_{A^c} = 1 - I_A$.

③ $I_{A \cap B} = I_A I_B$.

④ $I_{A \cup B} = I_A + I_B - I_A I_B$.

Fundamental Bridge Between Probability and Expectation

Theorem

There is a one-to-one correspondence between events and indicator r.v.s, and the probability of an event A is the expected value of its indicator r.v. I_A :

$$P(A) = E(I_A).$$

Example: Boole's Inequality

For any n events A_1, A_2, \dots, A_n ,

$$P\left(\bigcup_{i=1}^n A_i\right) \leq \sum_{i=1}^n P(A_i)$$

Solution: Boole's Inequality

Example: Inclusion-Exclusion Formula

For any events A_1, \dots, A_n :

$$P\left(\bigcup_{i=1}^n A_i\right) = \sum_i P(A_i) - \sum_{i < j} P(A_i \cap A_j) + \sum_{i < j < k} P(A_i \cap A_j \cap A_k) \\ - \dots + (-1)^{n+1} P(A_1 \cap \dots \cap A_n).$$

Solution: Inclusion-Exclusion Formula

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Moments of Indicator Methods

- Given n events A_1, \dots, A_n and indicators $I_j, j = 1, \dots, n$.
- $X = \sum_{j=1}^n I_j$: the number of events that occur
- $\binom{X}{2} = \sum_{i < j} I_i I_j$: the number of pairs of distinct events that occur
- $E\left(\binom{X}{2}\right) = \sum_{i < j} P(A_i \cap A_j)$
 - ▶ $E(X^2) = 2 \sum_{i < j} P(A_i \cap A_j) + E(X)$.
 - ▶ $\text{Var}(X) = 2 \sum_{i < j} P(A_i \cap A_j) + E(X) - (E(X))^2$.

Moments of Binomial Random Variables

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Poisson Distribution

Definition

An r.v. X has the *Poisson distribution* with parameter λ if the PMF of X is

$$P(X = k) = \frac{e^{-\lambda} \lambda^k}{k!}, \quad k = 0, 1, 2, \dots$$

We write this as $X \sim \text{Pois}(\lambda)$.

Example: Poisson Expectation & Variance

Poisson Approximation

Let A_1, A_2, \dots, A_n be events with $p_j = P(A_j)$, where n is large, the p_j are small, and the A_j are independent or weakly dependent. Let

$$X = \sum_{j=1}^n I(A_j)$$

count how many of the A_j occur. Then X is approximately $\text{Pois}(\lambda)$, with $\lambda = \sum_{j=1}^n p_j$.

Example: Birthday Problem Revisited

Poisson & Binomial

- Poisson \implies Binomial : **conditioning**
- Binomial \implies Poisson: **taking a limit**

Sum of Independent Poissons

Theorem

If $X \sim \text{Pois}(\lambda_1)$, $Y \sim \text{Pois}(\lambda_2)$, and X is independent of Y , then $X + Y \sim \text{Pois}(\lambda_1 + \lambda_2)$.

Poisson Given A Sum of Poissons

Theorem

If $X \sim \text{Pois}(\lambda_1)$, $Y \sim \text{Pois}(\lambda_2)$, and X is independent of Y , then the conditional distribution of X given $X + Y = n$ is $\text{Bin}(n, \lambda_1/(\lambda_1 + \lambda_2))$.

Poisson Approximation to Binomial

Theorem

If $X \sim \text{Bin}(n, p)$ and we let $n \rightarrow \infty$ and $p \rightarrow 0$ such that $\lambda = np$ remains fixed, then the PMF of X converges to the $\text{Pois}(\lambda)$ PMF. More generally, the same conclusion holds if $n \rightarrow \infty$ and $p \rightarrow 0$ in such a way that np converges to a constant λ .

Proof

Visitors to A Website

The owner of a certain website is studying the distribution of the number of visitors to the site. Every day, a million people independently decide whether to visit the site, with probability $p = 2 \times 10^{-6}$ of visiting. Give a good approximation for the probability of getting *at least three* visitors on a particular day.

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Typical Distance Measures

- Total Variation Distance
- Kullback–Leibler Divergence
- Jensen–Shannon Divergence
- Bhattacharyya Distance
- Wasserstein Distance (or called “Kantorovich–Rubinstein”)

Total Variation Distance

- Distance measure between two probability distributions
- Apply such measure to characterize the accuracy of Poisson approximation

Definition

The **total variation distance** between two distributions μ and ν on a countable set Ω is

$$\begin{aligned}d_{TV}(\mu, \nu) &= \| \mu - \nu \|_{TV} \\ &= \frac{1}{2} \sum_{x \in \Omega} |\mu(x) - \nu(x)|.\end{aligned}$$

Example

Let μ be the distribution with $\mu(1) = p$ and $\mu(0) = 1 - p$. Let ν be a Poisson distribution with mean p . Then we have $d_{TV}(\mu, \nu) \leq p^2$.

The Law of Small Numbers

Theorem

Given independent random variables Y_1, \dots, Y_n such that for any $1 \leq m \leq n$, $\mathbb{P}(Y_m = 1) = p_m$ and $\mathbb{P}(Y_m = 0) = 1 - p_m$. Let $S_n = Y_1 + \dots + Y_n$. Suppose

$$\sum_{m=1}^n p_m \rightarrow \lambda \in (0, \infty) \quad \text{as } n \rightarrow \infty,$$

and

$$\max_{1 \leq m \leq n} p_m \rightarrow 0 \quad \text{as } n \rightarrow \infty,$$

then

$$d_{TV}(S_n, \text{Poi}(\lambda)) \rightarrow 0 \quad \text{as } n \rightarrow \infty.$$

Gap of Poisson Approximation

- A bound on the gap due to Hodges and Le Cam (1960):

$$d_{TV}(S_n, Poi(\lambda)) \leq \sum_{m=1}^n p_m^2,$$

- by Stein-Chen method (C.Stein 1987) we can have a tighter bound on the gap:

$$d_{TV}(S_n, Poi(\lambda)) \leq \min(1, \frac{1}{\lambda}) \sum_{m=1}^n p_m^2.$$

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Probability Generating Function

Definition

The *probability generating function* (PGF) of a nonnegative integer-valued r.v. X with PMF $p_k = P(X = k)$ is the generating function of the PMF. By LOTUS, this is

$$E(t^X) = \sum_{k=0}^{\infty} p_k t^k.$$

The PGF converges to a value in $[-1, 1]$ for all t in $[-1, 1]$ since $\sum_{k=0}^{\infty} p_k = 1$ and $|p_k t^k| \leq p_k$ for $|t| \leq 1$.

Example: Generating Dice Probabilities

Let X be the total from rolling 6 fair dice, and let X_1, \dots, X_6 be the individual rolls. What is $P(X = 18)$?

Solution

PGF and Moments

Let X be a nonnegative integer-valued r.v. with PMF $p_k = P(X = k)$, and the PGF of X is $g(t) = \sum_{k=0}^{\infty} p_k t^k$, we have

- $E(X) = g'(t)|_{t=1}$
- $E(X(X - 1)) = g''(t)|_{t=1}$

PGF and Moments

PGF and Moments

Binomial PMF

Binomial Moments

Example: Pattern Matching

Suppose a coin with probability p for heads is tossed repeatedly, and we obtain a sequence of H and T (H denotes Head and T denotes Tail). Let N denote the number of toss to observe the first occurrence of the pattern “HH”. Find $E(N)$ and $\text{Var}(N)$.

Example: Pattern Matching

Example: Pattern Matching

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Probability Method

- Paul Erdős initiated this method: Erdős Method
- Widely used in information theory & combinatorics & theoretical computer science
- Main idea: to prove the existence of a structure with certain properties using probability or expectation

Principle I

- First we construct an appropriate probability space of structures.
- Then we show that a randomly chosen element in this space has the desired properties with positive probability

Theorem (The Possibility Principle)

Let A be the event that a randomly chosen object in a collection has a certain property. If $P(A) > 0$, then there exists an object with such property.

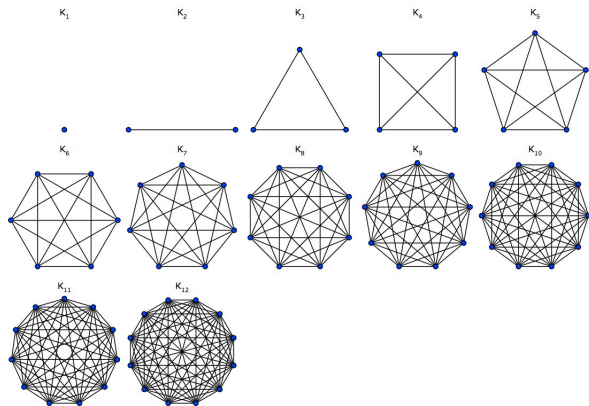
Principle II

Theorem (The Good Score Principle)

Let X be the score of a randomly chosen object. If $E(X) \geq c$, then there exists an object with a score of at least c .

Example: Graph Coloring

- Complete graph (clique): a simple undirected graph in which every pair of distinct vertices is connected by a unique edge.
- Complete graph K_n : a graph with n nodes and $\binom{n}{2}$ edges.



Example: Graph Coloring

Theorem

Given a complete graph K_n ($n \geq 3$), if $\binom{n}{m} 2^{-\binom{m}{2}+1} < 1$, then it is possible to color the edges of K_n with two colors so that it has no monochromatic K_m subgraph ($1 < m < n$).

Testing Polynomial Identities

- Randomized algorithms can be dramatically more efficient than their best known deterministic counterparts.
- Input two polynomials Q and R over n variables, with coefficients in some field, and decides whether $Q \equiv R$.
- Example: $Q(x_1, x_2) = (1 + x_1)(1 + x_2)$,
 $R(x_1, x_2) = 1 + x_1 + x_2 + x_1x_2$.
- n -variable polynomial $\prod_{i=1}^n (x_i + x_{i+1})$ expands into $O(2^n)$ monomials.

The Schwartz-Zippel Algorithm

- A Monte Carlo algorithm with a bounded probability of false positive and no false negative.
- Input polynomial $M(x_1, \dots, x_n)$ and test whether $M \equiv 0$ ($M = Q - R$).
- Assign values r_1, \dots, r_n chosen independently and uniformly at random from a finite set S to x_1, \dots, x_n .
- Test if $M(r_1, \dots, r_n) = 0$, outputting “Yes” if so and “No” otherwise.
- If “No”, then $M \not\equiv 0$.
- If “Yes”, it is possible that $M \not\equiv 0$ but r_1, \dots, r_n happens to be a zero of M .

Schwartz-Zippel Lemma

Lemma

Let $M \in F(x_1, x_2, \dots, x_n)$ be a non-zero polynomial of total degree $d \geq 0$ over a field F . Let S be a finite subset of F and let r_1, r_2, \dots, r_n be selected at random independently and uniformly from S . Then

$$P[M(r_1, r_2, \dots, r_n) = 0] \leq \frac{d}{|S|}.$$

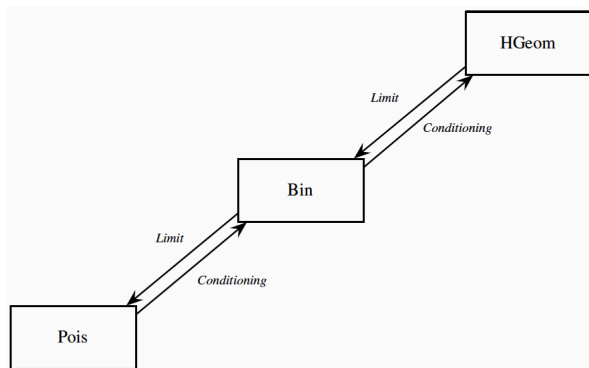
Remarks

- If we take the set S to have cardinality at least twice the degree of our polynomial ($|S| \geq 2d$), we can bound the probability of error (false positive) by $1/2$.
- This can be reduced to any desired small number by repeated trials.

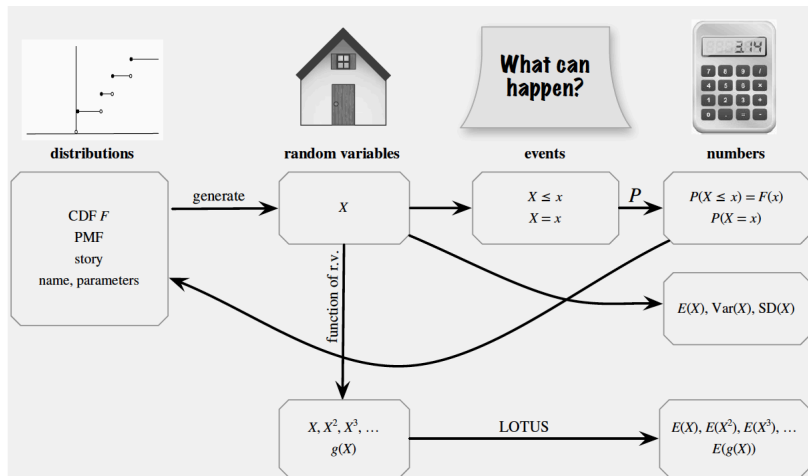
Summary 1

	With replacement	Without replacement
Fixed number of trials	Binomial	Hypergeometric
Fixed number of successes	Negative Binomial	Negative Hypergeometric

Summary 2



Summary 3



References

- Chapters 4 & 6 of **BH**
- Chapter 2 of **BT**